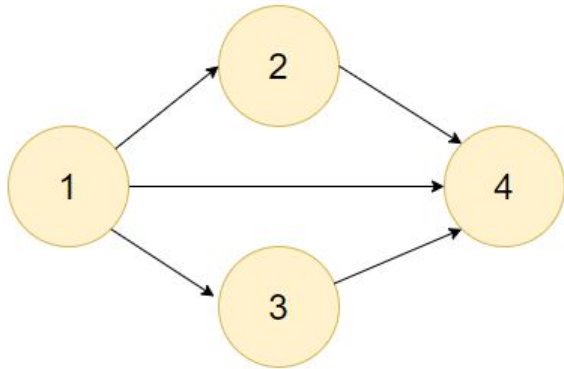


Algorithms: Graph Representation (Directed Graphs, Weighted Graphs)

Adjacency Matrix (Directed)

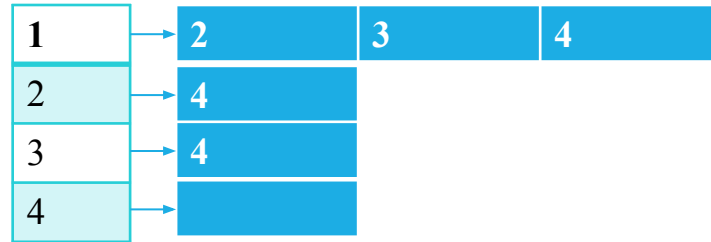
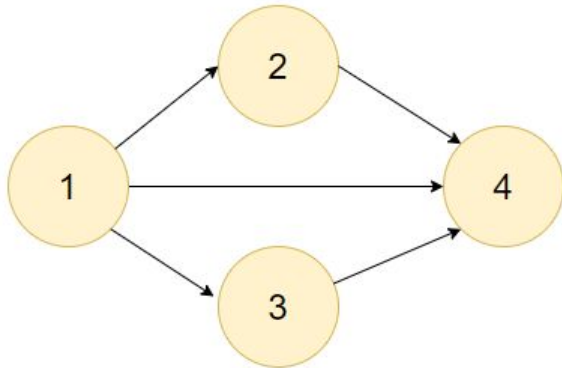


	1	2	3	4
1	0	1	1	1
2	0	0	0	1
3	0	0	0	1
4	0	0	0	0

Data Structure to Use: $(n * n)$ Matrix.

Example: `int adj [n][n];`

Adjacency List (Directed)



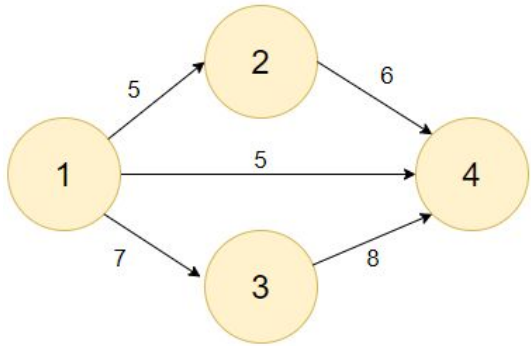
Data Structure to Use: 2D vector or array of linked list.

Example:

```
vector<int> adj [n];
```

```
vector< vector<int> > adj;
```

Adjacency Matrix (Directed and Weighted)

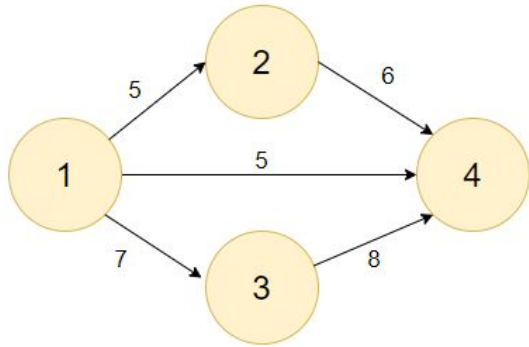


	1	2	3	4
1	0	5	7	5
2	0	0	0	6
3	0	0	0	8
4	0	0	0	0

Data Structure to Use: $(n * n)$ Matrix.

Example: `int adj [n][n];`

Adjacency List (Directed and Weighted)



Data Structure to Use:
2D vector or array of linked list.

Example:

```
vector< pair<int, int> > adj [n];
```

Adjacency Matrix vs Adjacency List

Adjacency Matrix:

- Space Complexity: $O(n^2)$
- Time complexity of Checking if an edge (u,v) exists: $O(1)$

- n : total number of nodes
- m : total number of edges
- d : degree (number edges connected to that node)

Adjacency List:

- Space Complexity: $O(n + m)$
- Time complexity of Checking if an edge (u, v) exists: $O(d)$

The benefit of using adjacency lists is that we can efficiently find the nodes to which we can move from a given node through an edge. For example, the following loop goes through all nodes to which we can move from node s :

```
for (auto u : adj[s]) {  
    // process node u  
}
```

